

Presto Training Series, Session 1: Using Advanced SQL Features In Presto

Try Presto: www.prestosql.io

David Phillips and Manfred Moser 29 July 2020

Today's Speakers



Manfred Moser
Developer, author,
and trainer at Starburst

Manfred is an open source developer and advocate. He is an Apache Maven committer, co-author of the book Presto: The Definitive Guide, and a seasoned trainer and conference presenter. He has trained over 20,000 developers for companies such as Walmart Labs, Sonatype, and Telus.



David Phillips
Co-creator of Presto and
CTO at Starburst

David is a co-creator of Presto, co-founder of the Presto Software Foundation, and CTO at Starburst. Prior to Starburst, David was a software engineer at Facebook, and held senior engineering positions at Proofpoint, Ning and Adknowledge.



Agenda

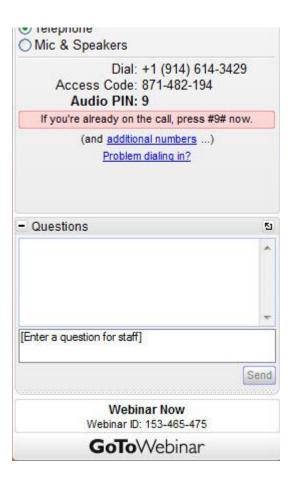
- Presto overview / review
- Using advanced SQL features in Presto
 - General SQL features
 - Using JSON
 - Advanced aggregation techniques
- Five minute break
 - Array and map functions
 - Window functions
- Q&A



Questions

Ask any time

- Use the meeting Questions feature
- Manfred screens, collects and interjects
- Dedicated Q&A in break and at end





Some advice for attendees

- This is a fast-paced overview don't try to follow along during class
- Instead focus and pay attention this is advanced SQL
- Use the demo video after class to setup Presto and CLI locally
- Learn at your own pace
- Use video recording and slides from class as reference to learn more
- Play with the TPC-H data sets
- Apply skills for your own use case



Presto Overview

... probably just a recap for you



What is Presto?



High performance ANSI SQL engine

- SQL support for any connected data source - SQL-on-anything
- Cost-based query optimizer
- Proven horizontal scalability



Separation of compute and storage

- Scale query processing and data sources independently
- Query storage directly
- No ETL or data integration necessary



Open source project

- Very active, large community
- User driven development
- Huge variety of users
- Prestosql.io



Presto everywhere

- No cloud vendor lock-in
- No storage engine vendor lock-in
- No Hadoop distro vendor lock-in
- No database lock in



Why use Presto?







Fastest time-to-insight

- High performance query processing
- Low barrier of entry for users
- Massive scalabilty
- High concurrency
- Direct access to storage

Lower cost

- Reduced need to copy and move data
- Avoid complex data processing
- Scale storage and compute independently
- Only run computes when processing queries
- One data consumption layer

Avoid data lock in

- No more data silos, departmental copies
- Query data with the existing skills and tools - SQL + BI tools
- Query any data source
- Move data
- Create optionality



General SQL Features



Format function

```
VALUES format('pi = %.5f', pi()),
       format('agent %03d', 7),
       format('$%,.2f', 1234567.89),
       format('%-7s,%7s', 'hello', 'world'),
       format('%2$s %3$s %1$s', 'a', 'b', 'c'),
       format('%1$tA, %1$tB %1$te, %1$tY', date '2006-07-04');
 pi = 3.14159
 agent 007
 $1,234,567.89
 hello , world
 b c a
 Tuesday, July 4, 2006
```



Simple case expressions

```
SELECT n,
       CASE n
           WHEN 1 THEN 'one'
           WHEN 2 THEN 'two'
           ELSE 'many'
       END AS name
FROM (VALUES 1, 2, 3, 4) AS t (n);
 n | name
     one
 2 | two
 3 | many
 4 | many
```



Searched case expressions

```
SELECT n,
       CASE
           WHEN n = 1 THEN 'aaa'
           WHEN n IN (2, 3) THEN 'bbb'
           ELSE 'ccc'
       END AS name
FROM (VALUES 1, 2, 3, 4) AS t (n);
n | name
 1 | aaa
 2 | bbb
 3 | bbb
 4 | ccc
```



IF expression



TRY expression

TRY avoids these failures:

```
Query failed: Division by zero

Query failed: Cannot cast 'abc' to INT

Query failed: integer addition overflow: 2000000000 + 2000000000
```



Lambda expressions overview

Lambda expression with one input:

$$x \rightarrow x + 8$$

Lambda expression with two inputs:

$$(x, y) \rightarrow x + y$$



Using JSON



JSON data type

```
SELECT
 json_parse('null') j_null,
 json_parse('true') bool,
 json_parse('"hello"') string,
 json_parse('123.45') number,
 json_parse('[1, 3, 5, 9]') array,
 json_parse('["hello", 123, {"xyz": 99, "abc": false}]') mixed;
j_null | bool | string | number | array |
                                                      mixed
------
null | true | "hello" | 123.45 | [1,3,5,9] | ["hello",123,{"abc":false,"xyz":99}]
```



Extraction using JSONPath



Casting from JSON



Partial casting from JSON

```
SELECT cast(v AS map(varchar, array(map(varchar, json))))
FROM (VALUES JSON '
    {"languages":
        [{"id": 123, "name": "Java", "data": [88,99]},
         {"id": 456, "name": "Python"}]}
') AS t (v);
                                   col0
 {languages=[{data=[88,99], name="Java", id=123}, {name="Python", id=456}]}
```



Formatting as JSON



Advanced Aggregation Techniques



Counting distinct items

How many unique customers do we have?

```
SELECT count(DISTINCT custkey) AS customers
FROM orders;

customers
------
99996
```

This gives an exact answer, but is slow and memory intensive.



Counting distinct items

Approximately how many unique customers do we have?

```
SELECT approx_distinct(custkey) AS customers
FROM orders;

customers
------
101655
```

This example has an error of 1.66%.



Counting distinct items

From the approx_distinct() documentation:

This function should produce a standard error of 2.3%, which is the standard deviation of the (approximately normal) error distribution over all possible sets. It does not guarantee an upper bound on the error for any specific input set.

The function uses the HyperLogLog algorithm to approximate the number of distinct items. The error depends on the unique values, not how many times they appear in the input. Both of these produce the same error:

- approx_distinct(x)
- approx_distinct(DISTINCT x)



Approximate percentiles

What is the order price at the 50th percentile?



Approximate percentiles

What are the order prices at the 10th, 20th, 50th, 90th, and 99th percentiles?



Associated max value

Find the clerk who has the most expensive order:



Associated max value using a row type

Find the clerk who has the most expensive order:



Associated max values

Find the clerks who have the most expensive orders:



Pivoting with conditional counting

Order counts by order priority, as separate columns:

```
SELECT
 count_if(orderpriority = '1-URGENT') AS urgent,
 count_if(orderpriority = '2-HIGH') AS high,
 count_if(orderpriority = '3-MEDIUM') AS medium,
 count_if(orderpriority = '4-NOT SPECIFIED') AS not_specified,
 count_if(orderpriority = '5-LOW') AS low
FROM orders;
urgent | high | medium | not_specified | low
300343 | 300091 | 298723 | 300254 | 300589
```



Pivoting with filtering

Order counts by order priority, as separate columns:

```
SELECT
 count(*) FILTER (WHERE orderpriority = '1-URGENT') AS urgent,
 count(*) FILTER (WHERE orderpriority = '2-HIGH') AS high,
 count(*) FILTER (WHERE orderpriority = '3-MEDIUM') AS medium,
 count(*) FILTER (WHERE orderpriority = '4-NOT SPECIFIED') AS not_specified,
 count(*) FILTER (WHERE orderpriority = '5-LOW') AS low
FROM orders;
 urgent | high | medium | not_specified | low
300343 | 300091 | 298723 | 300254 | 300589
```



Pivoting averages

Total order price by order priority, as separate columns:

```
SELECT
 avg(totalprice) FILTER (WHERE orderpriority = '1-URGENT') AS urgent,
 avg(totalprice) FILTER (WHERE orderpriority = '2-HIGH') AS high,
 avg(totalprice) FILTER (WHERE orderpriority = '3-MEDIUM') AS medium,
 avg(totalprice) FILTER (WHERE orderpriority = '4-NOT SPECIFIED') AS not_specified,
 avg(totalprice) FILTER (WHERE orderpriority = '5-LOW') AS low
FROM orders;
 urgent | high | medium | not_specified | low
151222.87 | 151553.28 | 151155.45 | 150792.44 | 151373.33
```



Aggregating a complex expression

What if we charge a premium based on order priority?

```
SELECT avg(totalprice *
           CASE
             WHEN orderpriority = '1-URGENT' THEN 1.10
             WHEN orderpriority = '2-HIGH' THEN 1.05
             ELSE 1.0
           END) / avg(totalprice) AS premium
FROM orders;
 premium
 1.03005
```



Aggregating into an array

Build an array from region names, in descending order:

```
SELECT array_agg(name ORDER BY name DESC) names
FROM region;
```

names

[MIDDLE EAST, EUROPE, ASIA, AMERICA, AFRICA]



Aggregating using a lambda

Compute the product of the values in the group:

```
SELECT name,
       reduce_agg(value, 1,
                 (a, b) -> a * b,
                 (a, b) -> a * b) AS product
FROM (VALUES ('x', 1), ('x', 3), ('x', 5),
            ('y', 2), ('y', 4), ('y', 6)) AS t (name, value)
GROUP BY name;
 name | product
x | 15
            48
```



Order-insensitive checksums

Compare data between tables by computing a checksum:



ROLLUP with single

```
SELECT orderpriority,
count(*) AS orders
FROM orders
GROUP BY ROLLUP(orderpriority)
ORDER BY orderpriority;
```

orderpriority	orders
	+
1-URGENT	300343
2-HIGH	300091
3-MEDIUM	298723
4-NOT SPECIFIED	300254
5-LOW	300589
NULL	1500000



ROLLUP with multiple

	returnflag		
F	+ А	-+·	1478493
F	N	Ī	38854
F	R	Ī	1478870
F	NULL	Ī	2996217
0	N	Ī	3004998
0	NULL	Ī	3004998
NULL	NULL	ī	6001215



CUBE

	returnflag	
	+	+
F	A	1478493
F	N	38854
F	R	1478870
F	NULL	2996217
0	N	3004998
0	NULL	3004998
NULL	A	1478493
NULL	N	3043852
NULL	R	1478870
NULL	NULL	6001215



GROUPING SETS

```
linestatus | returnflag |
                           items
           | A
                         1478493
                            38854
           l N
           l R
                         1478870
           NULL
                          2996217
           l N
                          3004998
           I NULL
                          3004998
NULL
           ΙΑ
                         1478493
NULL
           l N
                          3043852
NULL
                         1478870
                          6001215
NULL
           NULL
```



5 minute break

And if you stick around:

- Browse prestosql.io
- Join us on Slack
- Submit questions



Array and Map Functions



Creating arrays



Accessing array elements

Accessing an invalid subscript with [] fails:

```
Query failed: Array subscript must be less than or equal to array length: 8 > 5 Query failed: Array subscript is negative: -2
```



Sorting arrays

```
SELECT array_sort(ARRAY['a', 'xyz', 'bb', 'abc', 'z', 'b'],
    (x, y) \rightarrow CASE
                   WHEN length(x) < length(y) THEN -1
                   WHEN length(x) > length(y) THEN 1
                   ELSE 0
              END) AS sorted;
         sorted
 [a, z, b, bb, xyz, abc]
```



Matching elements

Do any, all, or none of the elements equal 8?



Filtering elements



Transforming elements



Converting arrays to strings

```
SELECT array_join(sequence(3, 7), '/') AS joined;
 joined
3/4/5/6/7
SELECT a,
      array_join(transform(a, e -> format('%,d', e)), ' / ') AS value
FROM (VALUES ARRAY[12345678, 987654321]) AS t (a);
            | value
 [12345678, 987654321] | 12,345,678 / 987,654,321
```



Computing array product



Unnesting an array

```
SELECT name
FROM (
    VALUES ARRAY['cat', 'dog', 'mouse']
) AS t (a)
CROSS JOIN UNNEST(a) AS x (name);
 name
 cat
 dog
 mouse
```



Unnesting an array with ordinality

```
SELECT id, name
FROM (
   VALUES ARRAY['cat', 'dog', 'mouse']
) AS t (a)
CROSS JOIN UNNEST(a) WITH ORDINALITY AS x (name, id);
 id | name
 1 | cat
  2 | dog
  3 | mouse
```



Creating maps

Create a map from arrays of keys and values:

```
SELECT map(ARRAY['x', 'y'], ARRAY[123, 456]);
```

Create a map from an array of entry rows:



Accessing map elements

Accessing an invalid key with [] fails:

```
Query failed: Key not present in map: bad
```



Unnesting a map

```
SELECT key, value
FROM (
    VALUES map_from_entries(ARRAY[('abc', 123), ('xyz', 456)])
) AS t (m)
CROSS JOIN UNNEST(m) AS x (key, value);
 key | value
 abc |
        123
 xyz
        456
```



Window Functions



Window function overview

Window functions run across rows of the result. Processing order:

- FROM and JOINs
- 2. WHERE
- 3. GROUP BY
- 4. HAVING
- 5. Window functions
- 6. SELECT
- DISTINCT
- 8. ORDER BY
- 9. LIMIT



Row numbering

Assign each region a unique number, in name order:

```
SELECT name,
      row_number() OVER (ORDER BY name) AS id
FROM region
ORDER BY name;
         | id
   name
AFRICA | 1
AMERICA | 2
ASIA
           1 3
EUROPE
MIDDLE EAST | 5
```



Row numbering order

Assign each region a unique number, in descending name order:

```
SELECT name,
      row_number() OVER (ORDER BY name DESC) AS id
FROM region
ORDER BY name;
        | id
   name
AFRICA | 5
AMERICA | 4
           | 3
ASIA
EUROPE
MIDDLE EAST | 1
```



Row numbering with limit

Assign each region a number, in descending name order, limited to three rows:

```
SELECT name,
       row_number() OVER (ORDER BY name DESC) AS row_number
FROM region
ORDER BY name
LIMIT 3;
             | id
    name
 AFRICA
 AMERICA
 ASIA
```



Rank

Assign a rank to each region, in descending name order:

name	 -	rank
AFRICA	+- 	 5
AMERICA	í	4
ASIA	Ī	3
EUROPE	Ī	2
MIDDLE EAST	Ī	1



Rank with ties

Assign a rank to each region, based on first letter of name:

name		rank
	+-	
AFRICA		1
AMERICA	I	1
ASIA	Ī	1
EUROPE	I	4
MIDDLE EAST	Ī	5



Dense rank with ties

Assign a rank to each region, based on first letter of name:

name	Ī	rank
	+-	
AFRICA		1
AMERICA	Ī	1
ASIA	Ī	1
EUROPE	Ī	2
MIDDLE EAST	Ī	3



Ranking without ordering

Assign a rank to each region:

```
SELECT name,
      rank() OVER (ORDER BY null) AS x,
      rank() OVER () AS y
FROM region
ORDER BY name;
         | x | y
   name
AFRICA | 1 | 1
 AMERICA
 ASIA
 EUROPE
MIDDLE EAST | 1 | 1
```



Row numbering without ordering

Assign a rank to each region:

```
SELECT name,
      row_number() OVER (ORDER BY null) AS x,
      row_number() OVER () AS y
FROM region
ORDER BY name;
        | x | y
   name
AFRICA | 1 | 1
 AMERICA | 2 | 2
 ASIA | 3 | 3
 EUROPE
       | 4 | 4
MIDDLE EAST | 5 | 5
```



Assigning rows to buckets

Assign rows into three buckets, in name order:

name	bucket
	+
AFRICA	1
AMERICA	1
ASIA	2
EUROPE	2
MIDDLE EAST	3



Percentage ranking

Percentage rank of rows, in name order:

name		percent
AFRICA	ï	0.0
AMERICA	i	0.25
ASIA	Ī	0.5
EUROPE	Ī	0.75
MIDDLE EAST	ī	1.0



Partitioning

Divide regions by first letter of name, then assign ranks:

name	I	rank
	+-	
AFRICA	1	1
AMERICA	1	2
ASIA	Ī	3
EUROPE	Ī	1
MIDDLE EAST		1



Partitioning on the same value

Assign a rank to each region:

```
SELECT name,
      rank() OVER (PARTITION BY null ORDER BY name) AS x,
      rank() OVER (ORDER BY name) AS y
FROM region
ORDER BY name;
        | x | y
   name
AFRICA | 1 | 1
 AMERICA | 2 | 2
 ASIA | 3 | 3
 EUROPE
       | 4 | 4
MIDDLE EAST | 5 | 5
```



Accessing leading and trailing rows

Access a value in the row behind and ahead of the current row:

```
SELECT name,
       lag(name) OVER (ORDER BY name) AS lag,
       lead(name) OVER (ORDER BY name) AS lead
FROM region
ORDER BY name;
                 lag
                            lead
   name
 AFRICA
             | NULL
                        | AMERICA
 AMERICA
             | AFRICA
                         ASIA
 ASIA
              | AMERICA |
                         EUROPE
 EUROPE
              | ASIA
                         MIDDLE EAST
 MIDDLE EAST | EUROPE
                       I NULL
```



Accessing leading and trailing rows

Access a value in the row behind and ahead of the current row, with default:

```
SELECT name,
       lag(name, 1, 'none') OVER (ORDER BY name) AS lag,
       lead(name, 1, 'none') OVER (ORDER BY name) AS lead
FROM region
ORDER BY name;
                 lag
                            lead
   name
 AFRICA
                       | AMERICA
             none
 AMERICA
             | AFRICA
                         ASIA
 ASIA
             | AMERICA |
                         EUROPE
 EUROPE
             | ASIA
                         MIDDLE EAST
 MIDDLE EAST | EUROPE
                       1 none
```



Accessing leading and trailing rows

Access a value two rows back and two rows ahead, with default:

```
SELECT name,
       lag(name, 2, 'none') OVER (ORDER BY name) AS lag2,
       lead(name, 2, 'none') OVER (ORDER BY name) AS lead2
FROM region
ORDER BY name;
                lag2 |
                          lead2
   name
 AFRICA
                       | ASIA
             none
 AMERICA
                        EUROPE
             none
 ASIA
             | AFRICA
                        MIDDLE EAST
 EUROPE
             | AMERICA | none
 MIDDLE EAST | ASIA
                       | none
```



Accessing leading and trailing rows with nulls

Access a value in the row behind and ahead of the current row, respecting nulls:

```
SELECT id, v,
      lag(v) OVER (ORDER BY id) AS lag.
      lead(v) OVER (ORDER BY id) AS lead
FROM (VALUES (1, 'a'), (2, 'b'), (3, null), (4, 'd'), (5, null)) AS t (id, v)
ORDER BY id:
id | v | lag | lead
----+----
 1 | a | NULL | b
 2 | b | a | NULL
 3 | NULL | b | d
 4 | d | NULL | NULL
 5 | NULL | d | NULL
```



Accessing leading and trailing rows without nulls

Access a value in the row behind and ahead of the current row, ignoring nulls:

```
SELECT id, x,
      lag(x) IGNORE NULLS OVER (ORDER BY id) AS lag,
      lead(x) IGNORE NULLS OVER (ORDER BY id) AS lead
FROM (VALUES (1, 'a'), (2, 'b'), (3, null), (4, 'd'), (5, null)) AS t (id, x)
ORDER BY id:
id | x | lag | lead
----+----
 1 | a | NULL | b
 2 | b | a | d
 3 | NULL | b | d
 4 | d | b | NULL
 5 | NULL | d | NULL
```



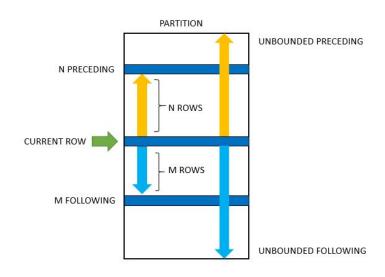
Window frames

Each row in a partition has a frame:

- ROWS: physical frame based on an exact number of rows
- RANGE: logical frame that includes all rows that are peers within the ordering

Examples:

- RANGE BETWEEN UNBOUNDED PRECEDING AND CURRENT ROW
- RANGE BETWEEN CURRENT ROW AND UNBOUNDED FOLLOWING
- ROWS BETWEEN 2 PRECEDING AND 2 FOLLOWING
- ROWS BETWEEN UNBOUNDED PRECEDING AND CURRENT ROW
- ROWS BETWEEN UNBOUNDED PRECEDING AND 5 FOLLOWING
- ROWS BETWEEN 3 PRECEDING AND UNBOUNDED FOLLOWING



Source: https://www.sqlitetutorial.net/sqlite-window-functions/sqlite-window-frame



Accessing the first value

```
SELECT name,
       first_value(name) OVER (
           ORDER BY name
           ROWS BETWEEN UNBOUNDED PRECEDING AND UNBOUNDED FOLLOWING
       ) AS value
FROM region;
             | value
   name
 AFRICA
             | AFRICA
 AMERICA
             | AFRICA
 ASIA
              | AFRICA
 EUROPE
              | AFRICA
 MIDDLE EAST | AFRICA
```



Accessing the last value

```
SELECT name,
       last_value(name) OVER (
           ORDER BY name
           ROWS BETWEEN UNBOUNDED PRECEDING AND UNBOUNDED FOLLOWING
       ) AS value
FROM region;
                value
   name
 AFRICA
             | MIDDLE EAST
 AMERICA
             | MIDDLE EAST
 ASIA
             | MIDDLE EAST
 EUROPE
             | MIDDLE EAST
 MIDDLE EAST | MIDDLE EAST
```



Accessing the Nth value

```
SELECT name,
      nth_value(name, 2) OVER (
          ORDER BY name
          ROWS BETWEEN UNBOUNDED PRECEDING AND UNBOUNDED FOLLOWING
       ) AS value
FROM region;
                 value
    name
 AFRICA
              | AMERICA
 AMERICA
              | AMERICA
 ASIA
              | AMERICA
 EUROPE
               AMERICA
 MIDDLE EAST
               AMERICA
```



Window frame ROWS vs RANGE

```
SELECT id, v,
      array_agg(v) OVER (ORDER BY id ROWS BETWEEN UNBOUNDED PRECEDING AND CURRENT ROW) AS rows,
      array_agg(v) OVER (ORDER BY id RANGE BETWEEN UNBOUNDED PRECEDING AND CURRENT ROW) AS range,
      array_agg(v) OVER (ROWS BETWEEN UNBOUNDED PRECEDING AND CURRENT ROW) AS rows_tie,
      array_agg(v) OVER (RANGE BETWEEN UNBOUNDED PRECEDING AND CURRENT ROW) AS range_tie
FROM (VALUES (1, 'a'), (2, 'b'), (3, 'c'), (3, 'd'), (5, 'e')) AS t (id, v);
id | v | rows | range | rows_tie | range_tie
 1 | a | [a] | [a] | [a, b, c, d, e]
 2 | b | [a, b] | [a, b] | [a, b, c, d, e]
 3 | c | [a, b, c] | [a, b, c, d] | [a, b, c] | [a, b, c, d, e]
 3 | d | [a, b, c, d] | [a, b, c, d] | [a, b, c, d] | [a, b, c, d, e]
 5 | e | [a, b, c, d, e] | [a, b, c, d, e] | [a, b, c, d, e] | [a, b, c, d, e]
```



Rolling and total sum

```
SELECT V,
      sum(v) OVER (ORDER BY v ROWS BETWEEN UNBOUNDED PRECEDING AND CURRENT ROW) AS rolling,
     sum(v) OVER () total
FROM (VALUES 1, 2, 3, 4, 5) AS t (v);
v | rolling | total
1 | 1 | 15
2 | 3 | 15
3 |
    6 |
             15
4 | 10 | 15
5 |
     15 | 15
```



Partition sum

```
SELECT p, v,
      sum(v) OVER (
          PARTITION BY p ORDER BY v
          ROWS BETWEEN UNBOUNDED PRECEDING AND CURRENT ROW) AS sum
FROM (VALUES ('a', 1), ('a', 2), ('a', 3), ('b', 4), ('b', 5), ('b', 6)) AS t (p, v);
p | v | sum
a | 2 | 3
a | 3 | 6
b | 5 | 9
b | 6 | 15
```



Wrapping up



Presto Training Series

Join the Presto creators again for more:

- Understanding and Tuning Query Processing with Martin (12 Aug)
- Securing Presto with Dain (26 Aug)
- Configuring and Tuning Presto Performance with Dain (9 Sept)



Presto Summit series

Diverse information about Presto and real world usage

- State of Presto <u>recording available</u>
- Presto as Query Layer at Zuora <u>recording available</u>
- Presto Migration at Arm Treasure Data recording available
- Presto for Analytics at Pinterest <u>19 Aug</u>



And finally ...

- Learn more from our website and documentation at <u>prestosql.io</u>
- Join us on slack at <u>prestosql.io/slack</u>
- Get a free digital copy of <u>Presto: The Definitive Guide</u>
- Thank you for hanging out with us
- See you next time



Your question Our answers ...

